

OSIsoft Overview

Developments in the PI System
Family of Products

Gregg Le Blanc
PI System Product Manager

Why Scalability?

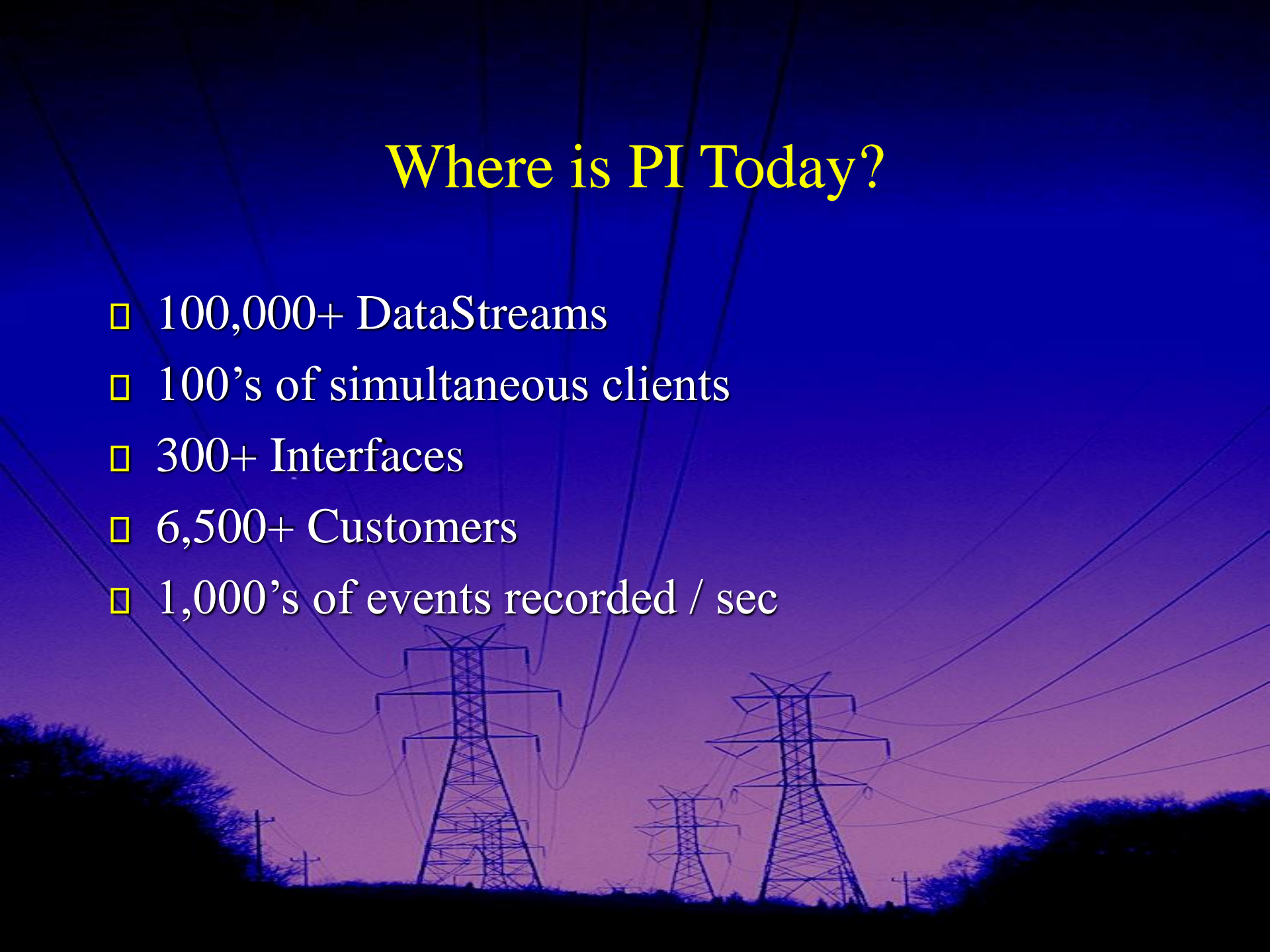
- More data can be tracked every day
 - Phones have GPS location data
 - Real-time monitoring in the home
 - Electrical usage
 - Security systems
 - Fire / Flood detection
 - IT Applications
 - Every machine on a network needs monitoring
- Example:
 - $700,000 \text{ homes} \times 10 \text{ data points} = 7,000,000 \text{ points}$

Benefits

- IT Monitoring can detect
 - Intrusions
 - Impending hardware problems
 - Diagnose software problems
 - Virus traffic
- Example:
 - Because WiredCity could detect and record the amount of inbound traffic from the Nimda worm, they got a refund from MCI Worldcom

Where is PI Today?

- ❑ 100,000+ DataStreams
- ❑ 100's of simultaneous clients
- ❑ 300+ Interfaces
- ❑ 6,500+ Customers
- ❑ 1,000's of events recorded / sec



Issues Within Scalability

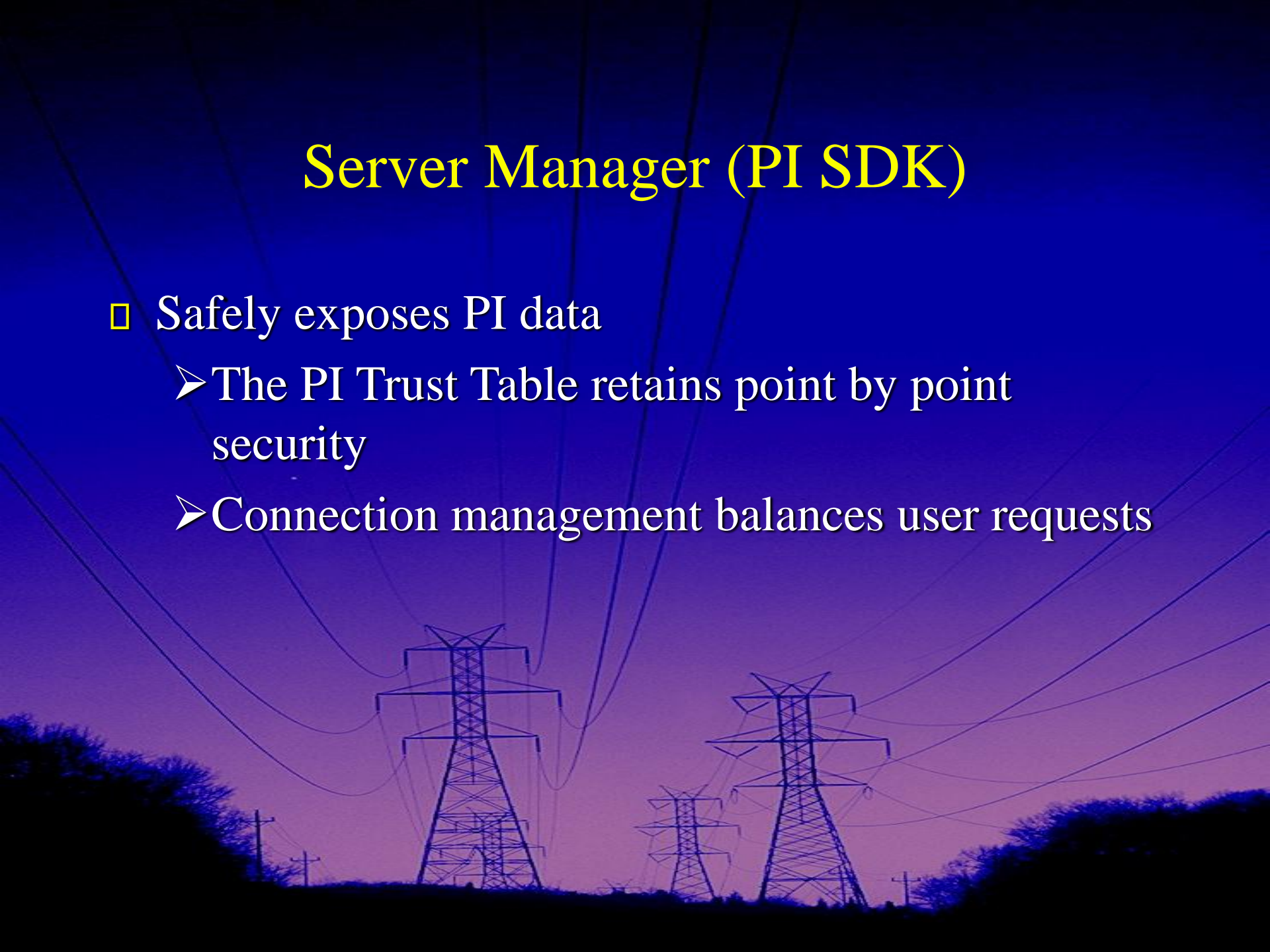
- A scalable system has several aspects
 - Users served
 - Calculations and analysis
 - Integration issues
 - Data capacity
 - Data throughput
 - Security
 - Infrastructure

User Load & Integration

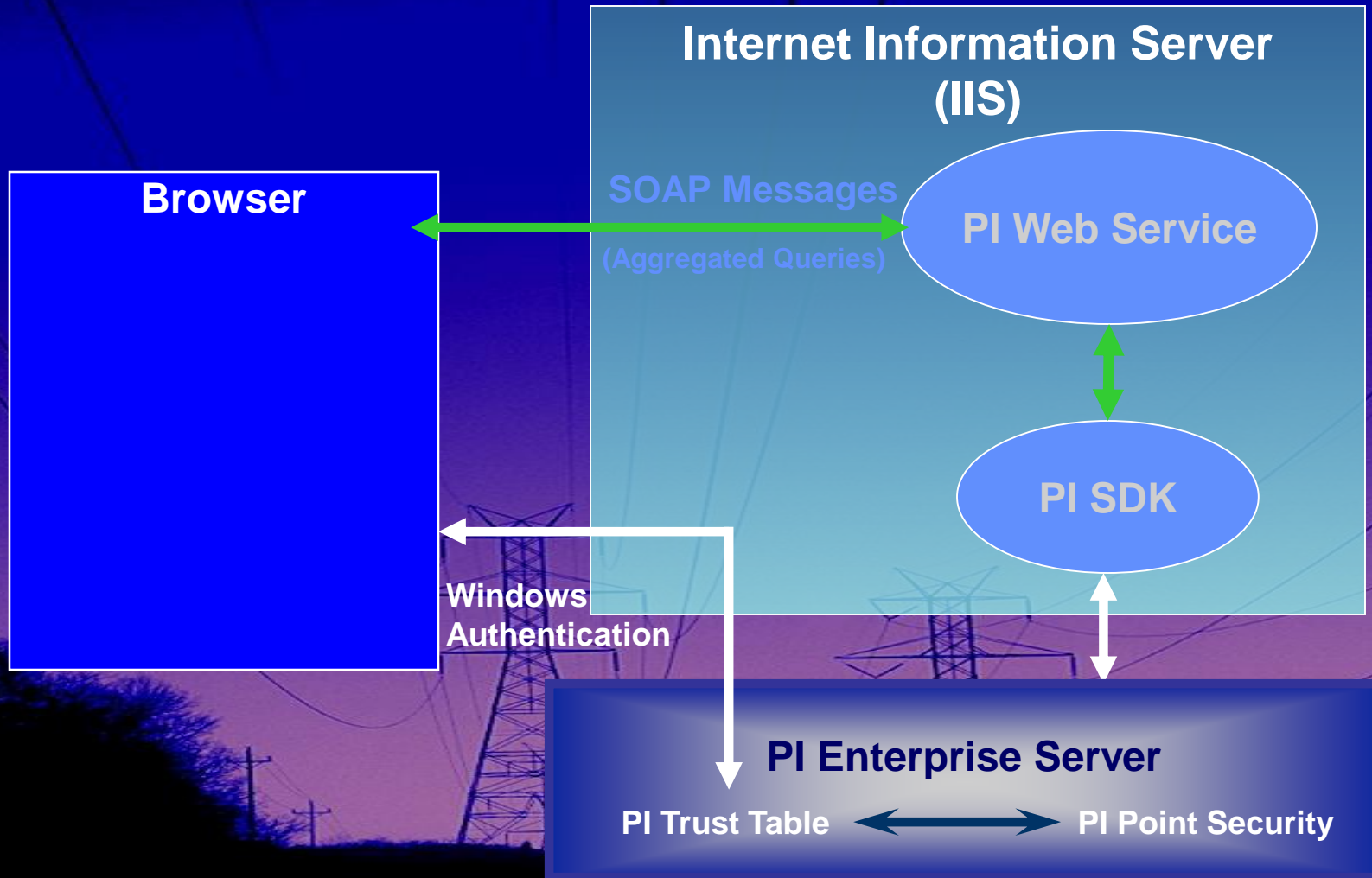
- Users need fast access to the data
 - Client / Server users
 - Stateless Users
- Client / Server users connect directly to PI
- Stateless Users
 - Broader audience
 - PI ICE
 - Web services and integration

Server Manager (PI SDK)

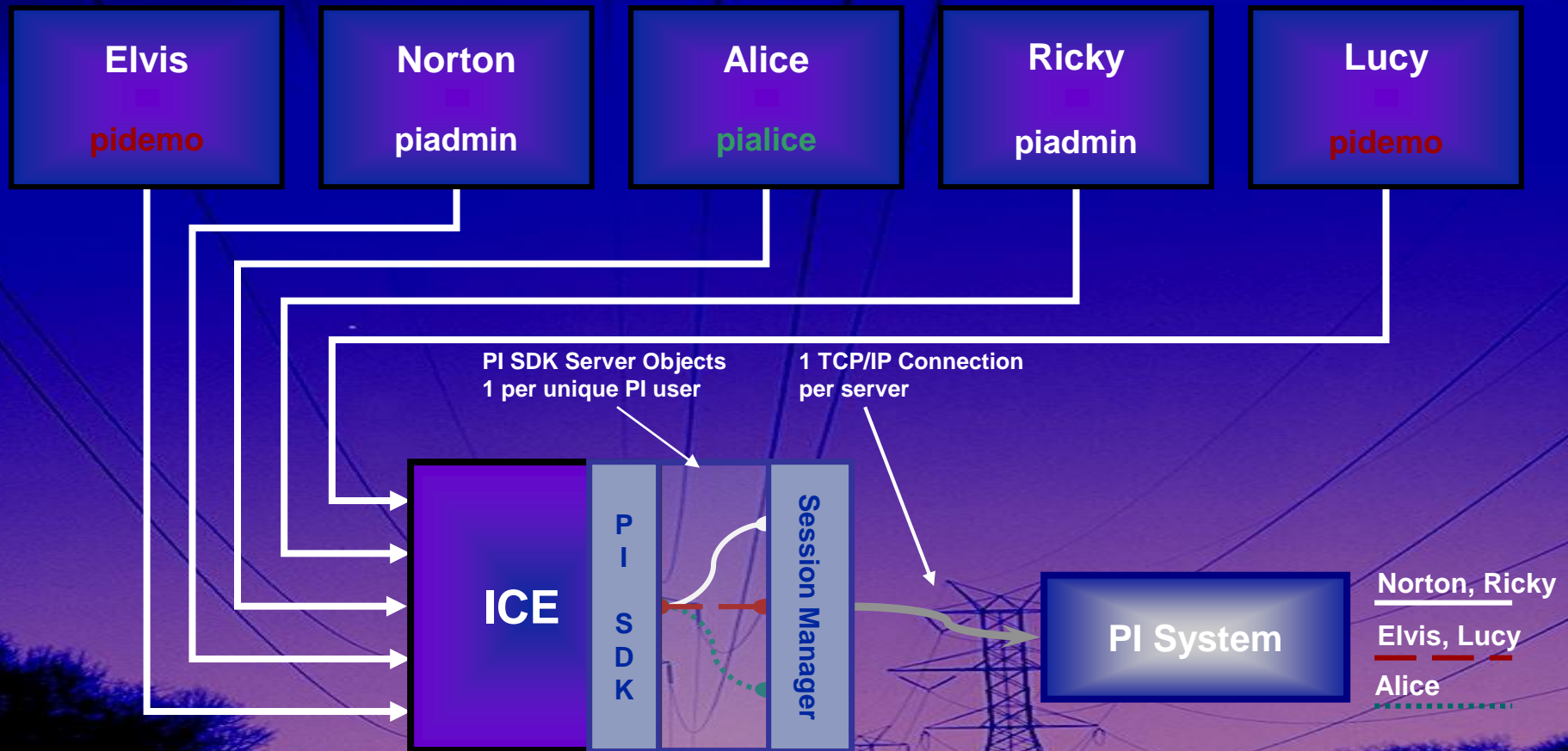
- Safely exposes PI data
 - The PI Trust Table retains point by point security
 - Connection management balances user requests



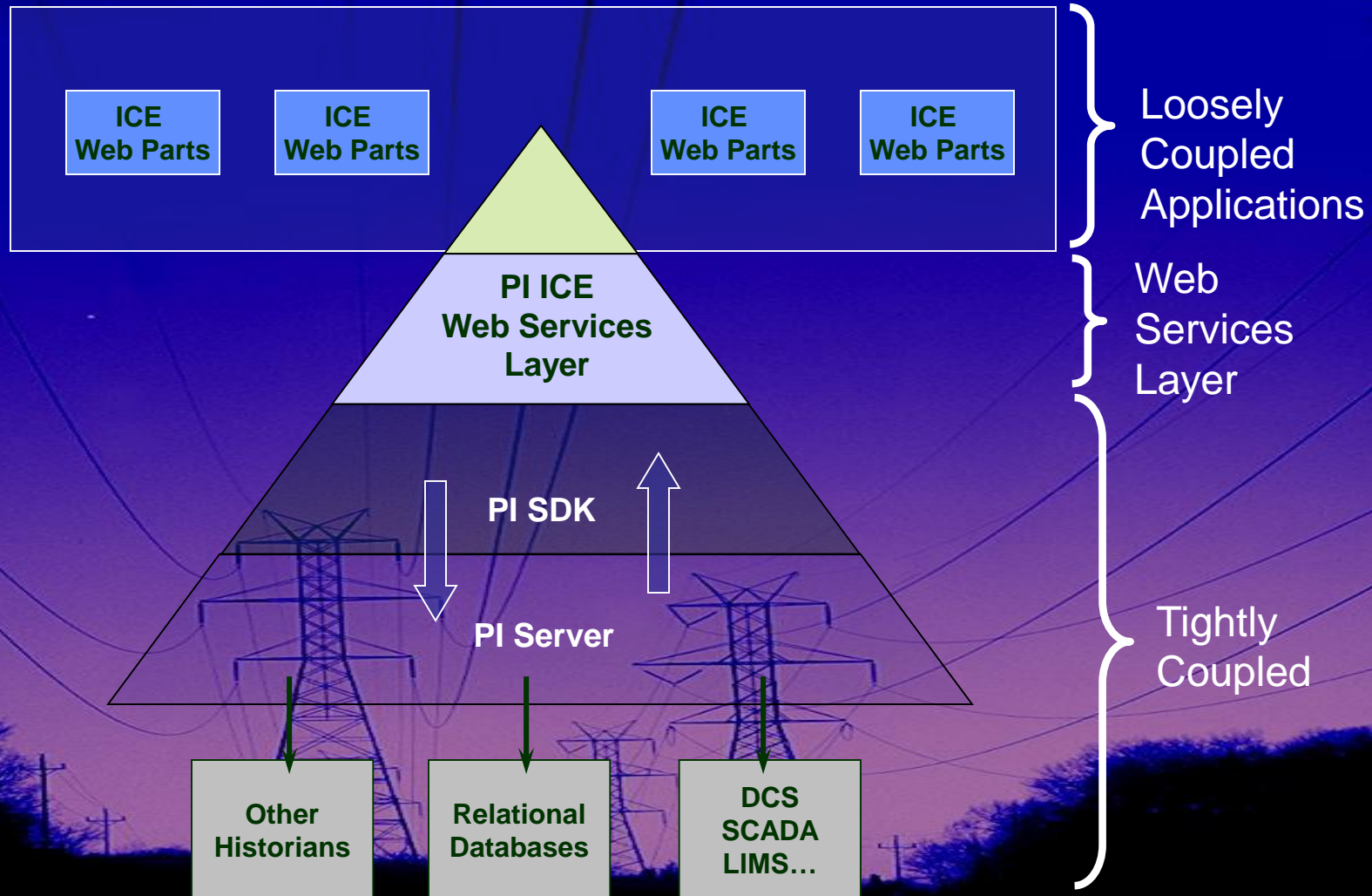
Security in ICE



Connections in ICE



Underlying Web Service Structure



Loosely Coupled Computing

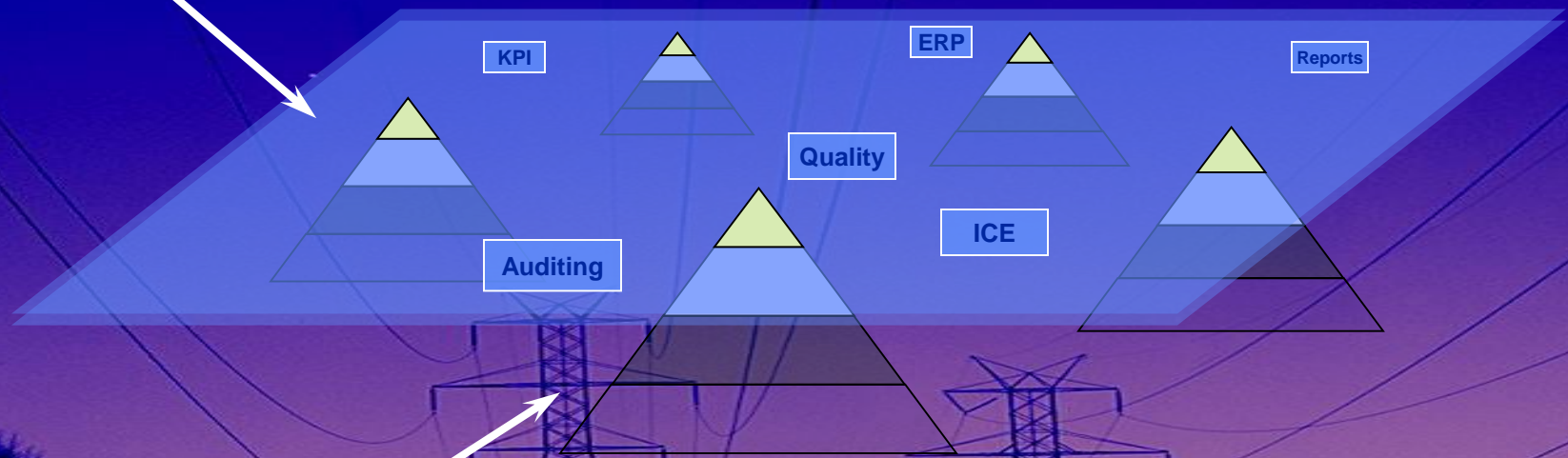
- The idea of Web services
- Allows for computing in an environment that is:
 - Asynchronous
 - Stateless
 - Platform independent
 - Geographically independent
- Will replace some “tightly coupled” computing

Tightly Coupled Computing

- Infrastructure-level ties between systems
 - Uses API calls
 - Usually proprietary
 - Not as flexible
- Examples:
 - PI Interfaces to other information systems
 - OLEDB / ODBC connectors
 - Point-to-Point, Middleware, and other integration

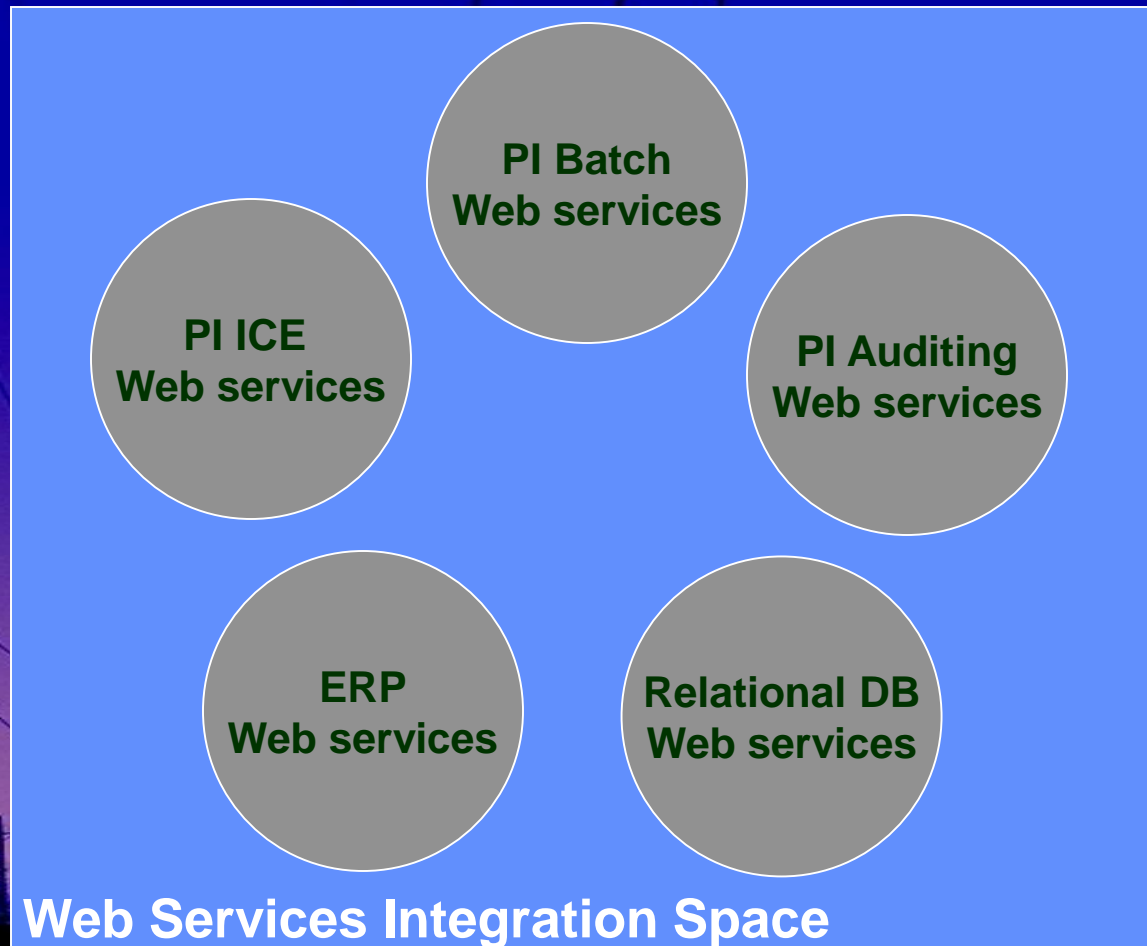
Web Service Integration

Loosely coupled applications built on
Web Service Integration Space



Traditional Tightly Coupled Integration Space

Web Services



Infrastructure

- ❑ OS/soft is moving to .NET
- ❑ Older applications will migrate
- ❑ New applications will leverage the .NET Framework
- ❑ Management tools will use .NET
 - New tools coming for SMT (Systems Management Tools)
- ❑ Parts of the PI System will be accessible through .NET

Distributed Analysis

- Server Based
 - Performance Equations
- Desktop Based
 - Spreadsheets based on DataLink
- Distributed Calculations
 - PI ACE
- Server load reduced
- Calculations can be reused

Security Concerns

□ Isn't Microsoft a vulnerability?

➤ Over the last year

- 26 vulnerabilities in Apache
- 22 in Microsoft's IIS

➤ Key issue: system maintenance

□ What about Nimda?

➤ 250,000 systems in 9 hours

➤ At least \$2.4B in damages

➤ Infection started July 13, 2001

➤ Patch was available June 18, 2001 (MS01-033)

Security Response

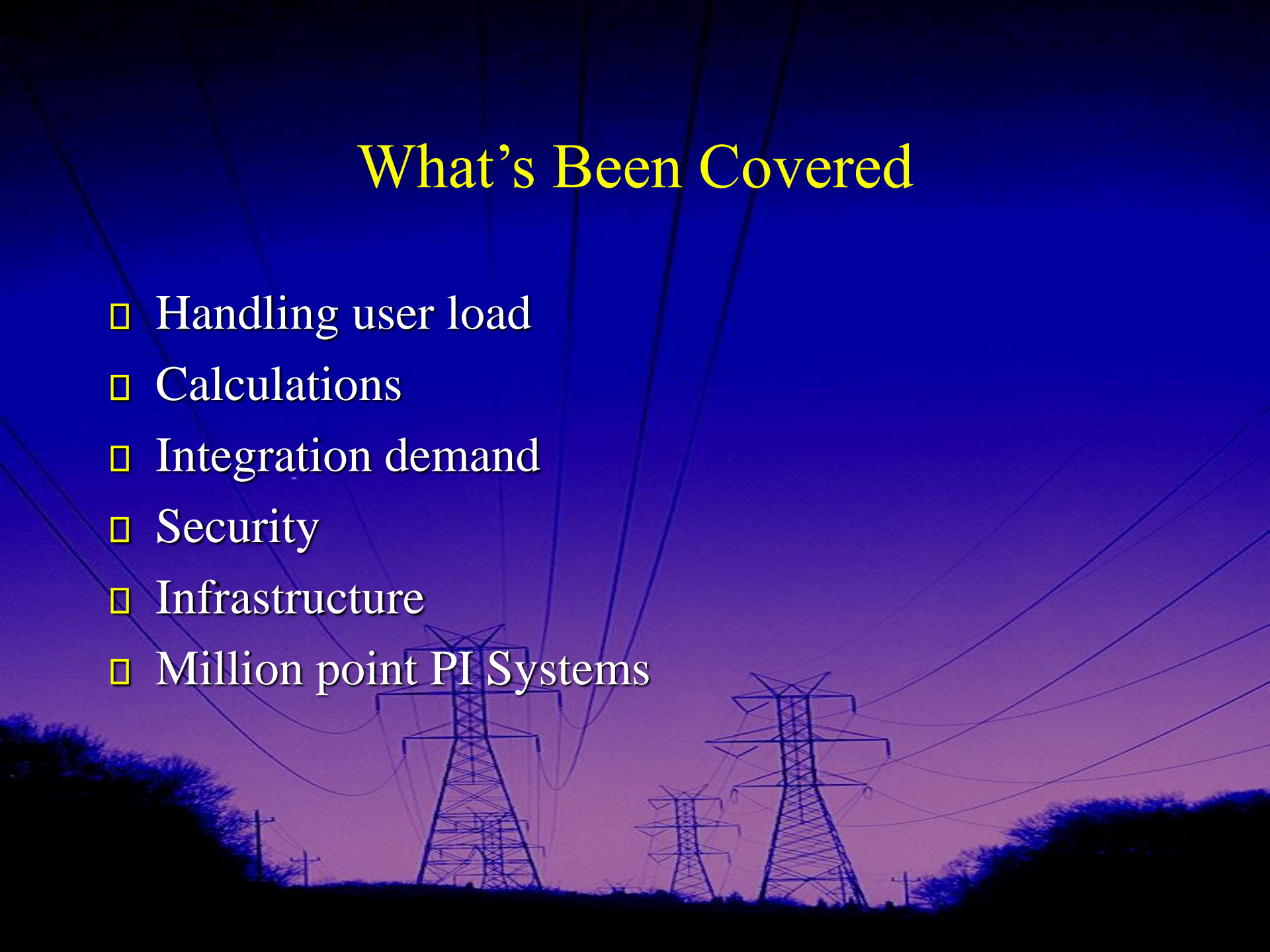
- ❑ More accountability: Microsoft
- ❑ NIPC (National Infrastructure Protection Center) works with NERC and Microsoft to develop security procedures and standards
- ❑ Key precautions still remain
 - Guard against social engineering
 - System maintenance
 - Unknown vulnerabilities
 - ❑ WiFi, wardriving
 - ❑ Unprotected modem lines used for support

OSIsoft Security

- ❑ PI Point Security
- ❑ Trust Table
- ❑ Single direction PI to PI transfers
- ❑ Auditing
 - A new database that records changes to PI
 - ❑ Points
 - ❑ Values
 - ❑ Module changes

What's Been Covered

- Handling user load
- Calculations
- Integration demand
- Security
- Infrastructure
- Million point PI Systems



History of the Historian

- PI 1
 - 1983 HP-1000
 - 1985 Vax-VMS
 - 1988 PINet
 - Client/Server Architecture
- PI 2
 - 1992 Vax-VMS
 - 1994 Alpha-OpenVMS
- PI 3
 - 1993 Design Work

History of the Historian

□ PI 3

➤ Big 4 Unix

- HP-UX

- IBM-AIX

- Dec-OSF/1

- Sun-Solaris

➤ Windows NT 3.51

➤ Development Language C++

History of the Historian

□ PI 3

➤ Key concepts of PI 2 were used

- Snapshot
- Compression
- Archive cache
- Archive navigation
- PINet

History of the Historian

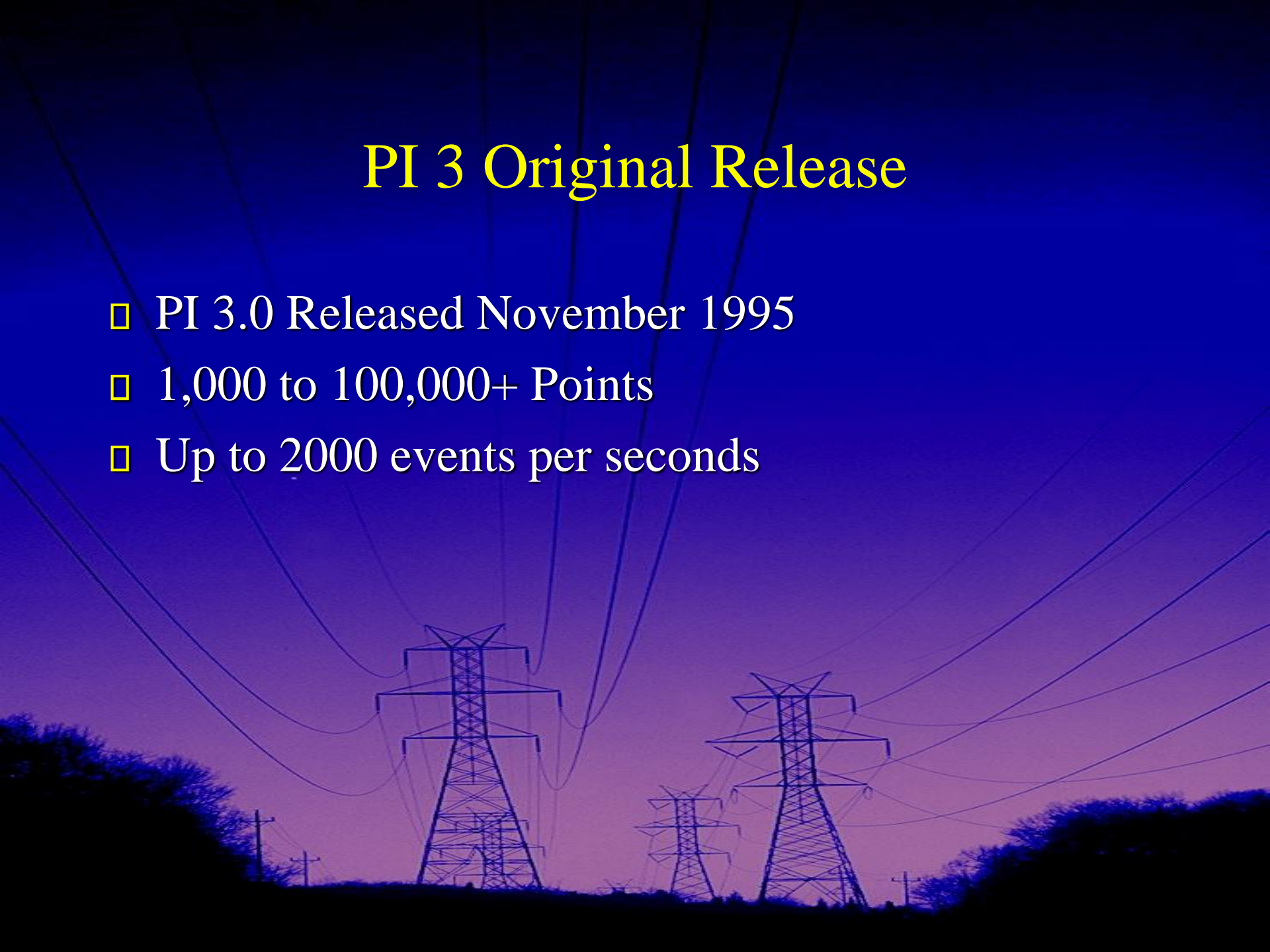
□ PI 3

➤ Key new concepts

- Multi-process
- RPC Based Inter-process communication
 - PI SDK
- Historization of many different data types
 - Doubles
 - Strings
 - BLOBs
- Common code base for all platforms
 - 64 bit support

PI 3 Original Release

- ❑ PI 3.0 Released November 1995
- ❑ 1,000 to 100,000+ Points
- ❑ Up to 2000 events per seconds



PI 3.0 to PI 3.3

□ How much can you do in 7 years?

➤ Concentrated on features

- Alarm
- ACE
- Totalizer
- PI SDK
- Batch Database
- Module Database
- Audit
- NT Security
- COM Connectors
- Development Infra-structure
 - Automated builds and testing
- Bug fixes
 - (just a few)

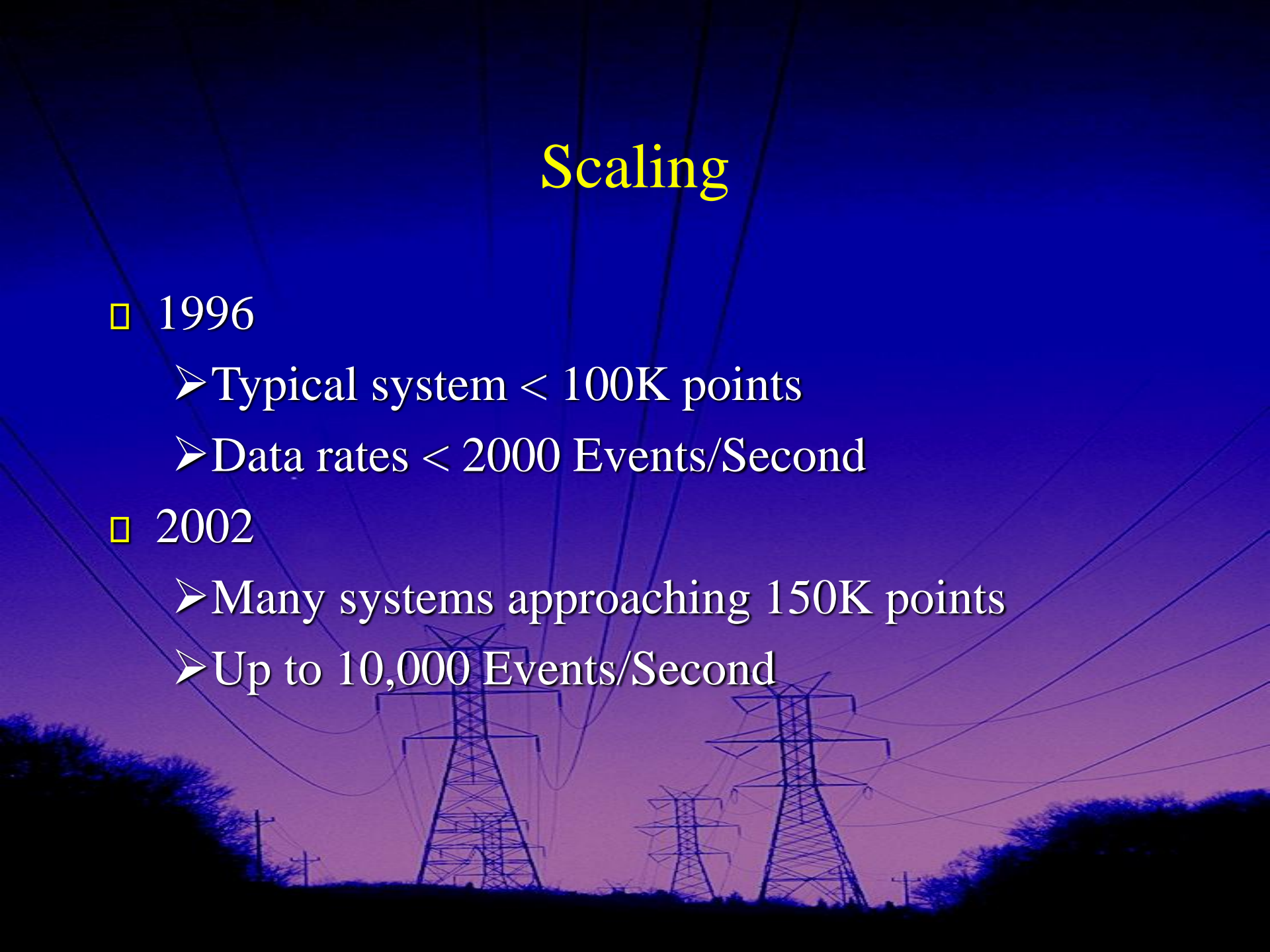
Scaling

□ 1996

- Typical system < 100K points
- Data rates < 2000 Events/Second

□ 2002

- Many systems approaching 150K points
- Up to 10,000 Events/Second



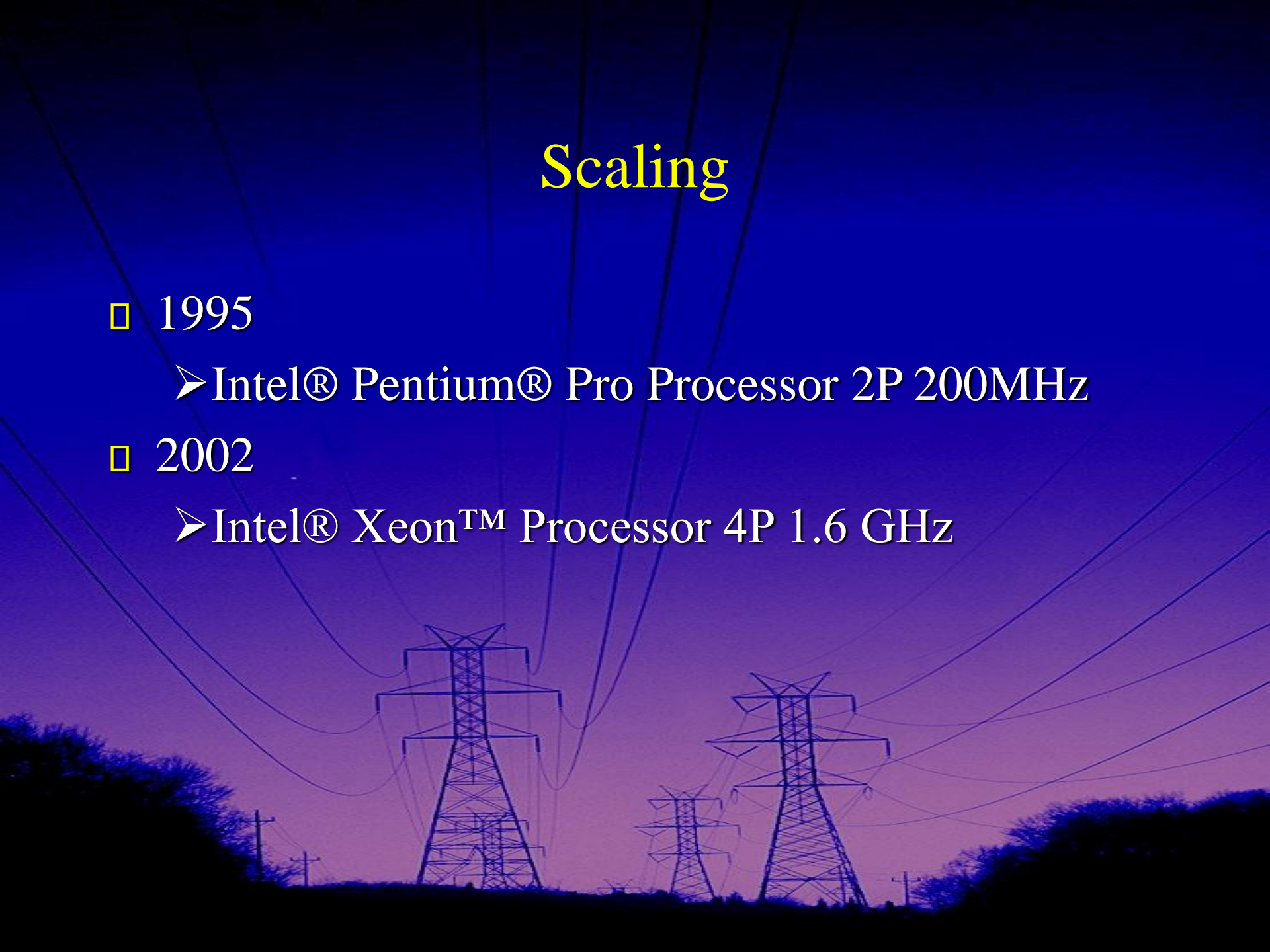
Scaling

□ 1995

➤ Intel® Pentium® Pro Processor 2P 200MHz

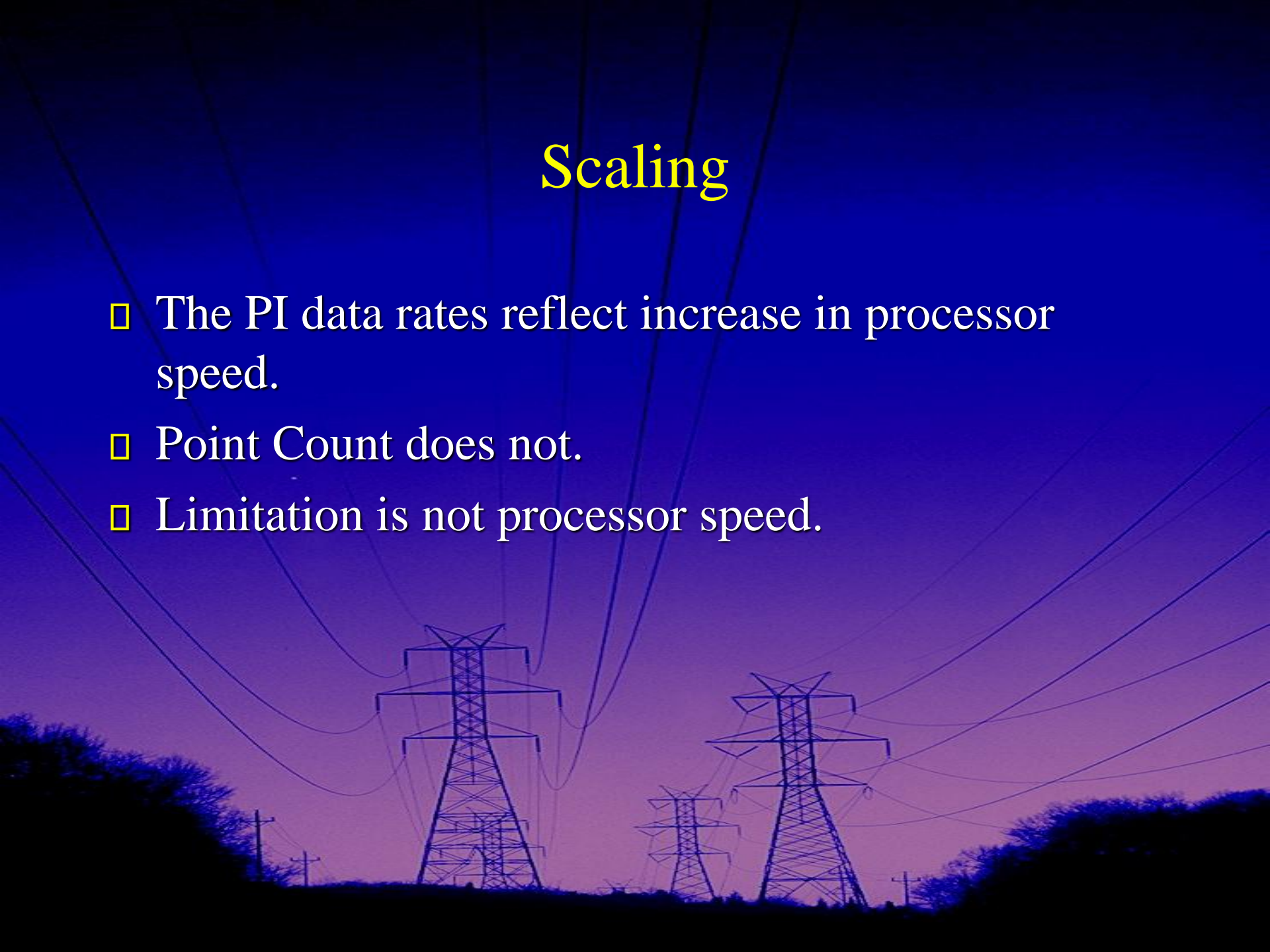
□ 2002

➤ Intel® Xeon™ Processor 4P 1.6 GHz



Scaling

- ❑ The PI data rates reflect increase in processor speed.
- ❑ Point Count does not.
- ❑ Limitation is not processor speed.



Limitations in PI 3.3

- Memory

- 2 GB of virtual memory per process

- Archive cache

- Inter-process communication

- Compressed events from snapshot to archive

- Serialization of RPCs

- Archive sub-system handles one call at a time

Memory

- PI Databases are memory resident
 - Point database
 - Module database
 - Snapshot
- PI Archive Cache

PI 3.3 Cache Memory Issues

- ❑ Cache record can be quite large.
- ❑ Adding a single event requires entire cache record
- ❑ Systems receiving data for most points will have a large memory footprint
 - Lots of cache activity
 - ❑ Pushing records out of memory
 - ❑ Reading records into memory

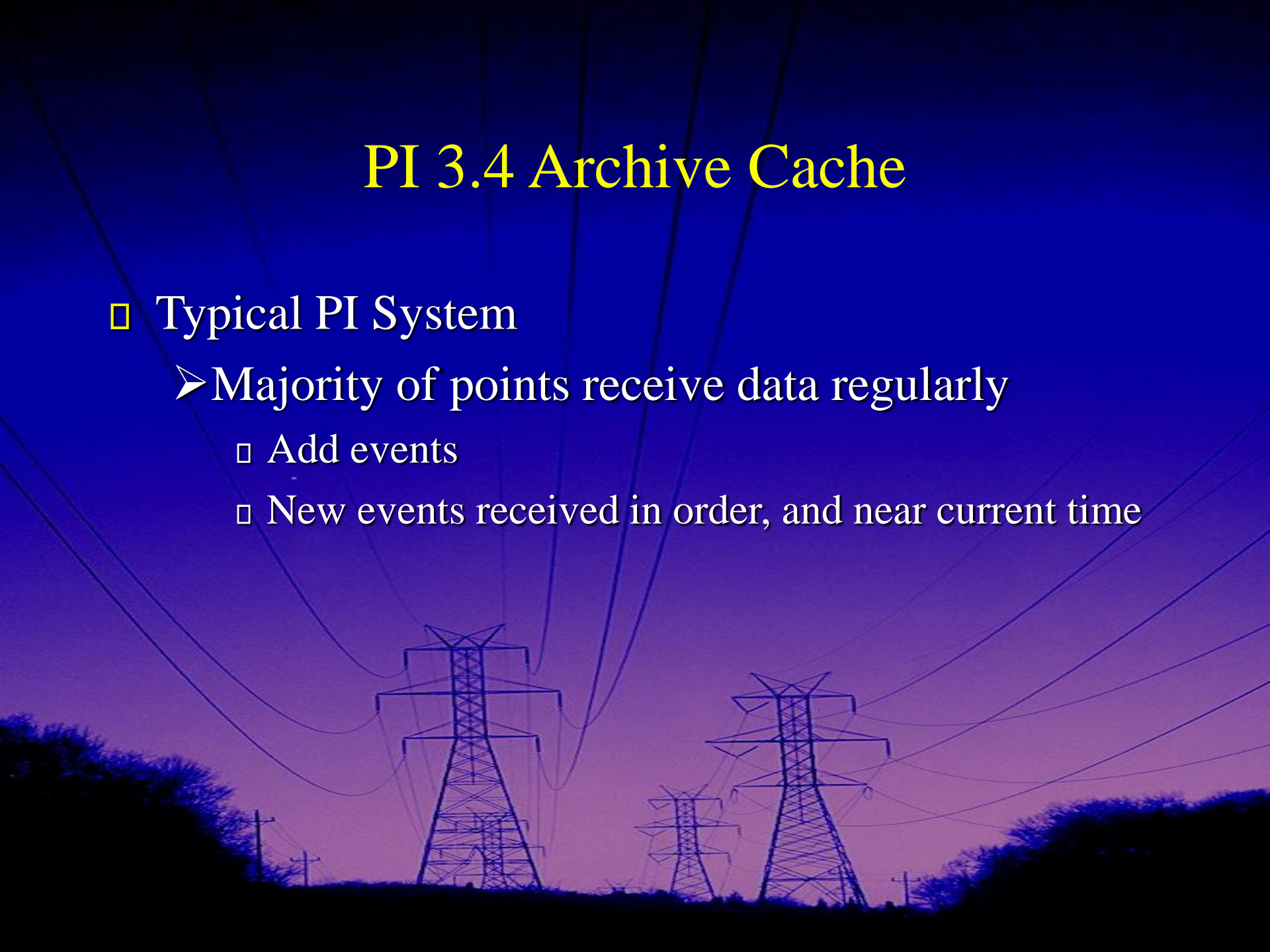
PI 3.4 Archive Cache

- Typical PI System

- Majority of points receive data regularly

- Add events

- New events received in order, and near current time



The PI 3.4 Cache

- Addresses Memory Issues
 - Smaller memory footprint
 - More efficient handling of new events
 - Same efficiency reading data
- Design Tradeoffs
 - More complicated

3.3 Snapshot – Archive Inter-process Communication

□ Current Mechanism

- Event received is processed by the compression algorithm.
- If passes compression event is added to in-memory event queue
- Events in queue are packaged and sent via an RPC to the archive process

3.3 Snapshot Buffering

- ❑ Archive unable to process compressed events
 - Backups
 - Extremely high data rates
 - System problems
- ❑ Snapshot must write to “event queue” file
- ❑ Archive must read from the Snapshot event queue
- ❑ Result
 - 2 physical disk I/O's happen on same file

3.4 Snapshot – Archive

Inter-process Communication

- Changed physical file to memory mapped file
 - File system feature
- Mapped by Snapshot and Archive
- Snapshot writes compressed events to “file”
- Archive reads compressed events from “file”
 - Synchronization techniques required to prevent corruption
- Buffering case handled by file system

Limitation 3: RPC Serialization

- ❑ PI 3.3 all sub-systems are single threaded
- ❑ Requests to sub-systems are serialized
 - Archive example—3 simultaneous calls
 - ❑ Archive Summary
 - ❑ Plot values
 - ❑ Compressed values



RPC Serialization Solution

- ❑ Multiple Threads
 - Threads are the obvious approach
 - Unix and NT have very powerful and easy to use threading models
- ❑ Conceptually Simple
- ❑ But, easy to write bugs

Multi-Threading in PI 3.4

- PI 3 had some threading since the original release
 - PI Net Manager
 - Sub-systems use a thread for reading messages
- PI 3.4 Sub-system Level thread model
 - Main “house keeping” thread
 - Read thread
 - Message pump
 - Pool of worker threads

Multi-Threading in PI 3.4

- PI 3.4 Sub-system Level thread model
 - Every sub-system is multi-threaded
 - Sub-system must implement locks for optimal performance
 - Worker thread pool is configurable
 - Runtime
 - Kill threads
 - Suspend threads
 - Change priority
 - Add worker threads
 - Delete worker threads

Multi-Threading in PI 3.4

- What does this get us?
 - The archive call from Won't kill your system.
 - What will 10 archive calls form do to my system?

```
Long time1 = 100  
Long time2 = 1134289  
piar_summary ( 1023453, &time1, &time2, &rval, &pctgood, ARCTOTAL )
```

PI on 64 Bit Windows

- One advantage: larger addressable memory
 - 32 bit: 2 GB
 - 64 bit: 16 – 128 GB
 - XP-64: 16 GB
 - .Net Enterprise Server: 64 GB
 - .Net Data Center: 128 GB

PI on 64 Bit Windows

- ❑ Disadvantages:

- Slower

- Very sensitive to properly optimized code

- ❑ Servers only

- 64 bit Windows is not intended to replace 32 bit desktops.

- Only makes sense to port server applications

- ❑ 64 bit development will be done on 32 bit machines

PI on 64 Bit Windows

- PI 3 is designed for 64 bit
 - Supports Dec OSF/1
 - Original release of PI supported Dec's 64 bit Unix and still does—HP Tru64 Unix.
- PI 3.4 will be 64 bit ready
 - Minor differences between Windows and Unix 64 bit
 - 3.4 will address these minor issues
 - 3.4 will compile cleanly

Summary

- Point Count limited by archive cache memory
- 3.4 archive cache is significantly smaller and more efficient
- Memory mapped files for configuration data
 - Point Database
- Tested million point system on typical 32 bit server class machine
 - 2P 1 GHz Pentium 4 / 4 GB Memory
 - Several million points possible
- 64 bit, perhaps 50 Million points?

Summary

- Snapshot – Archive inter-process communication
 - Significantly faster the RPC mechanism
 - Memory mapped file shared by both processes
 - Backed by disk to insure no data loss
- RPC Serialization
 - Eliminated by threading
 - Process several calls simultaneously

Questions?

